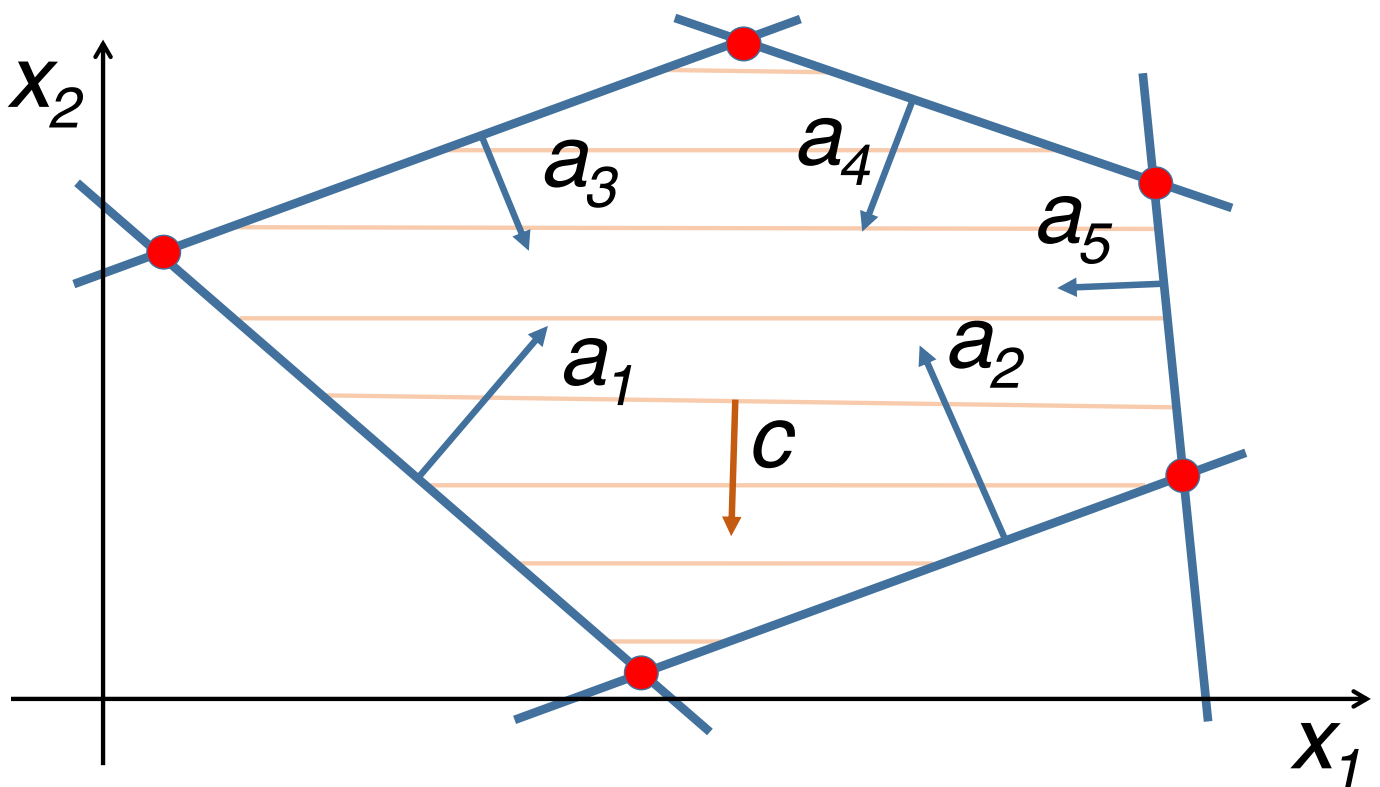


# Introduction to Linear Programming

## What is LP

Generally speaking, all problems with linear objective and linear equalities/inequalities constraints could be considered as Linear Programming. However, there are some widely accepted formulations.

$$\begin{aligned} \min_{x \in \mathbb{R}^n} c^\top x & \quad (\text{LP.Basic}) \\ \text{s.t. } Ax \leq b \end{aligned}$$



for some vectors  $c \in \mathbb{R}^n$ ,  $b \in \mathbb{R}^m$  and matrix  $A \in \mathbb{R}^{m \times n}$ . Where the inequalities are interpreted component-wise.

## Standard form

This form seems to be the most intuitive and geometric in terms of visualization. Let us have vectors  $c \in \mathbb{R}^n$ ,  $b \in \mathbb{R}^m$  and matrix  $A \in \mathbb{R}^{m \times n}$ .

$$\begin{aligned} \min_{x \in \mathbb{R}^n} c^\top x & \quad (\text{LP.Standard}) \\ \text{s.t. } Ax = b \\ x_i \geq 0, i = 1, \dots, n \end{aligned}$$

## Canonical form

$$\begin{aligned} & \min_{x \in \mathbb{R}^n} c^\top x && \text{(LP. Canonical)} \\ \text{s.t. } & Ax \leq b \\ & x_i \geq 0, \quad i = 1, \dots, n \end{aligned}$$

## Real world problems

### Diet problem

Imagine, that you have to construct a diet plan from some set of products: 🍌 🍰 🍗 🥚 🐟 . Each of the products has its own vector of nutrients. Thus, all the food information could be processed through the matrix  $W$ . Let also assume, that we have the vector of requirements for each of nutrients  $r \in \mathbb{R}^n$ . We need to find the cheapest configuration of the diet, which meets all the requirements:

$$\begin{aligned} & \min_{x \in \mathbb{R}^p} c^\top x \\ \text{s.t. } & Wx \geq r \\ & x_i \geq 0, \quad i = 1, \dots, n \end{aligned}$$



$$W \in \mathbb{R}^{n \times p},$$

Requirements

$$r \in \mathbb{R}^n$$

Proteins  
Carbs  
Fats  
Calories  
Vitamin D

$c \in \mathbb{R}^p$  - cost per 100 g

$$\begin{aligned} & \min_{x \in \mathbb{R}^p} c^\top x \\ & Wx \geq r \end{aligned}$$

## How to retrieve LP

## Basic transformations

Inequality to equality by increasing the dimension of the problem by  $m$ .

$$Ax \leq b \leftrightarrow \begin{cases} Ax + z = b \\ z \geq 0 \end{cases}$$

unsigned variables to nonnegative variables.

$$x \leftrightarrow \begin{cases} x = x_+ - x_- \\ x_+ \geq 0 \\ x_- \geq 0 \end{cases}$$

## Chebyshev approximation problem

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_\infty \leftrightarrow \min_{x \in \mathbb{R}^n} \max_i |a_i^\top x - b_i|$$

$$\begin{aligned} & \min_{t \in \mathbb{R}, x \in \mathbb{R}^n} t \\ & \text{s.t. } a_i^\top x - b_i \leq t, \quad i = 1, \dots, n \\ & \quad -a_i^\top x + b_i \leq t, \quad i = 1, \dots, n \end{aligned}$$

## $l_1$ approximation problem

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_1 \leftrightarrow \min_{x \in \mathbb{R}^n} \sum_{i=1}^n |a_i^\top x - b_i|$$

$$\begin{aligned} & \min_{t \in \mathbb{R}^n, x \in \mathbb{R}^n} \mathbf{1}^\top t \\ & \text{s.t. } a_i^\top x - b_i \leq t_i, \quad i = 1, \dots, n \\ & \quad -a_i^\top x + b_i \leq t_i, \quad i = 1, \dots, n \end{aligned}$$

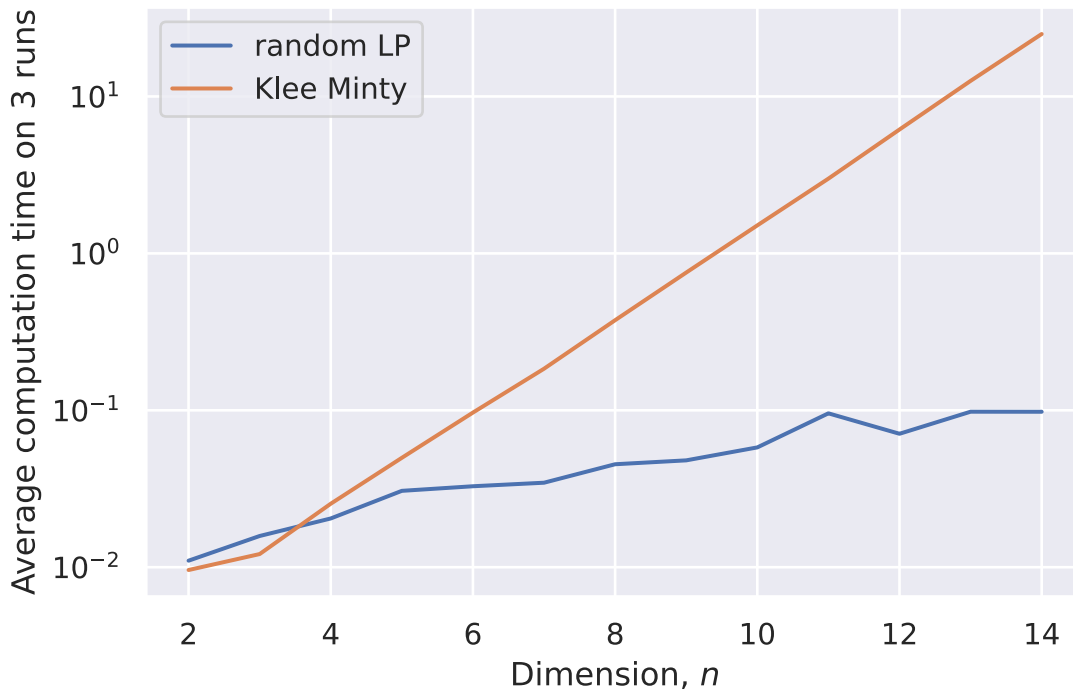
## Idea of simplex algorithm

# Convergence

## Klee Minty example

In the following problem simplex algorithm needs to check  $2^n - 1$  vertexes with  $x_0 = 0$ .

$$\begin{aligned} & \max_{x \in \mathbb{R}^n} 2^{n-1}x_1 + 2^{n-2}x_2 + \dots + 2x_{n-1} + x_n \\ \text{s.t. } & x_1 \leq 5 \\ & 4x_1 + x_2 \leq 25 \\ & 8x_1 + 4x_2 + x_3 \leq 125 \\ & \dots \\ & 2^n x_1 + 2^{n-1}x_2 + 2^{n-2}x_3 + \dots + x_n \leq 5^n \quad x \geq 0 \end{aligned}$$



## Summary

- A wide variety of applications could be formulated as the linear programming.
- Simplex algorithm is simple, but could work exponentially long.
- Khachiyan's ellipsoid method is the first to be proved running at polynomial complexity for LPs. However, it is usually slower than simplex in real problems.
- Interior point methods are the last word in this area. However, good implementations of simplex-based methods and interior point methods are similar for routine applications of linear programming.

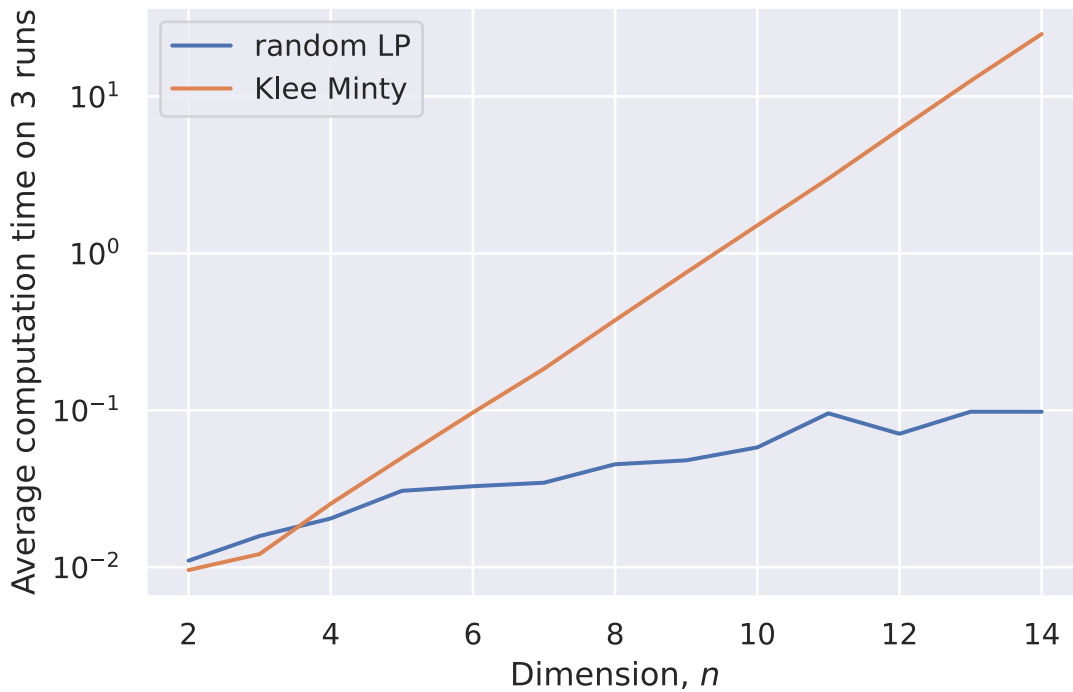
## Code

 [Open in Colab](#)

## Materials

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